# Device Boot-up Security Restoration with Post-Quantum Split-key KEM

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\***Akman et al**., "Split keys for station-to-station (STS) protocols", Journal of surveillance security and safety, vol. 4, no.3, 2023.

#### **Premises of the Work**

Embedded devices boot-up security relies on non-PQC cryptographic primitives, so it cannot be expect to be secure in the future.
A naïve way of employing non-PQC protocols along with PQC protocols is not a secure solution (Mosca's Theorem).
A solution is a protocol for SecureBoot security restoration by a firmware update.

#### **Attacker Model and Protocol Requirements**

#### **SecureBoot**

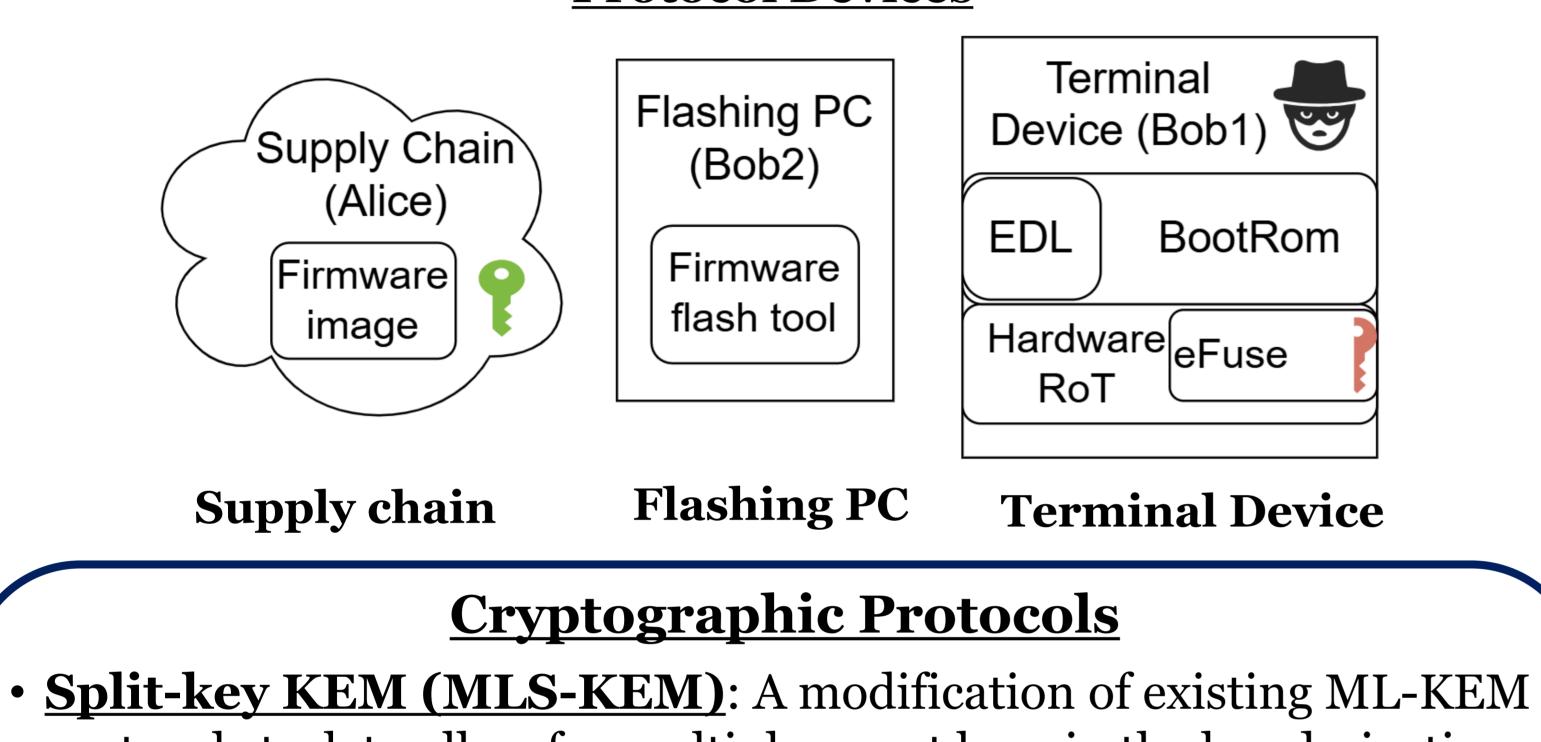
• SecureBoot is a verification mechanism for the software run during boot. It checks software signatures and checksums, which ensures that they are trusted and executed in correct order.

#### **One run of the MLS-KEM Protocol**

Alie $pk_B$ ,		<b>Bob 2</b> pk <sub>A</sub> , sk <sub>B2</sub>	<b>Bob 1</b> <i>pk</i> <sub>A</sub> , <i>sk</i> <sub>B1</sub>
<b>1</b> . ( <b>C</b> , <b>K</b> ) = MLS-KEM.Encaps $(pk_B)$			
	2. C		

3. (K2, v') = MLS-KEM.Decaps( $\mathbf{C}, sk_{B2}$ )

The attacker can (i) break legacy crypto, (ii) infect the system with malicious code and (iii) read secret keys already present on the device.
The protocol requirements are then (i) the protocol must be based on PQC primitives, (ii) it must be run in BootROM in Emergency Download Mode (EDL), and (iii) the protocol must not rely only on the keys/secrets on the compromised device



#### **Protocol Devices**

#### 4. (C, K2)

5. (K1, v) = MLS-KEM.Decaps(C,  $sk_{B1}$ )

6.  $K = MLS-KEM.Reconst(K1, K2, sk_{B1}, v)$ 

#### **MLS-KEM**

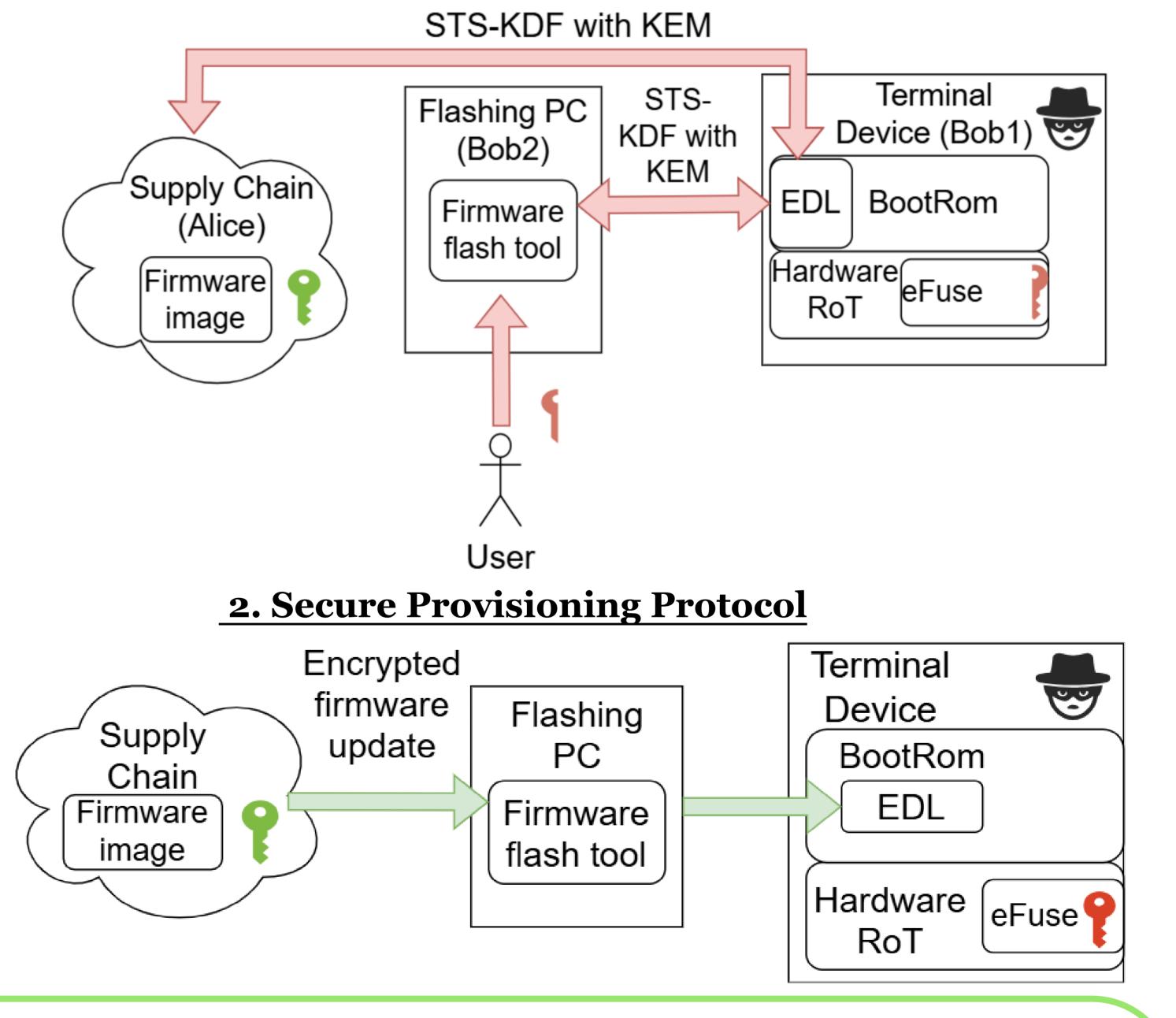
•MLS-KEM is a modified ML-KEM (NIST standardized version of CRYSTALS-Kyber) lattice-based KEM with split-keys.

• **MLS-KEM** performance is comparable to ML-KEM provided that similar parameters are used.

## **Protocol Architecture**

• The solution to boot-up security restoration: use two protocols for **secure provisioning**, and for **endpoint attestation**.





protocol stack to allow for multiple secret keys in the key derivation process.

- **STS-KDF with split-key KEM\*:** A protocol between three parties, two of which have **split keys**, and the result is a **session key**. Assuming the split-key KEM is secure, STS-KDF is formally proven to provide (i) **confidentiality** of the session key, (ii) **mutual authentication**, (iii) **forward security**.
- If split-key KEM is QC secure  $\Rightarrow$  STS-KDF is PQC.

#### The Use Case

- The protocol is between a **user** (with a split-key), a **supply chain** (with a public key) and the infected **terminal device** (with a split-key).
- Basic assumptions, (i) user's key half has not leaked, (ii) the terminal device is compromised (key half has leaked), (iii) BootRom is not compromised

# **The Protocol Functionality**

• To download the firmware update to the compromised device, the protocol executes the **endpoint attestation protocol** and then the

## secure provisioning protocol.

- 1. First, the **User** initialises the **Terminal Device** in **EDL** model. This ensures that no compromised software is run before the firmware update.
- 2. The **User** provides his keyshare to initialise the **endpoint attestation protocol** (STS-KDF with split-KEM). The result is the establishment of an encrypted connection between the **Terminal Device** and the **Supply Chain**.
- 3. The **secure provisioning protocol** is executed, and the result is a **firmware update** to the **Terminal Device**. The **Terminal Device** then installs the firmware update in **EDL** mode, restoring the boot-up security of the device.

# Helsinki System Security Lab (HSSL)

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